



HUMAN CAPITAL
NATIONAL COHESION STRATEGY



EUROPEAN UNION
EUROPEAN
SOCIAL FUND



Philosophy and Methodology of Sciences

Tomasz Placek

Truth and proof

(lecture 4)

Project co-financed by the European Union under the European Social Fund

DEFINITION: THEOREM OF CPL

A wff α is a CPL theorem iff there is a finite sequence of wffs $\gamma_1, \dots, \gamma_k$ (proof), such that its last element is α ($\gamma_k = \alpha$) and each element is

either (1) a CPL axiom,

or (2) is derived from the preceding elements by the rule MP.

Illustration

$p \rightarrow p$ CPL' s law of identity

1. $(p \rightarrow (q \rightarrow q)) \rightarrow (p \rightarrow p)$

2. $p \rightarrow (q \rightarrow q)$

3. $p \rightarrow p$

Ax. 1 $\alpha/p \rightarrow (q \rightarrow q), \beta/p$

Ax. 1 $\alpha/p, \beta/q$

MP 1, 2

DEFINITION: provability in CPL

A wff α is provable in CPL from the set Φ of wffs ($\Phi \vdash_p \alpha$) iff there is a finite sequence of wff's $\gamma_1, \dots, \gamma_k$, such that its last element is α ($\gamma_k = \alpha$) and each element of the sequence is

either (1) **an element of Φ (a premise)**,

or (2) an axiom of CPL,

or (3) is derived from preceding element by the rule MP

Provability (syntactic consequence) in CPL

$$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$$

1. α

premise

2. β

premise

$$3. (\overset{\alpha}{\alpha} \rightarrow \overset{\beta}{\alpha}) \rightarrow ((\overset{\alpha}{\alpha} \rightarrow \overset{\gamma}{\beta}) \rightarrow (\overset{\alpha}{\alpha} \rightarrow (\overset{\beta}{\alpha} \wedge \overset{\gamma}{\beta})))$$

Ax. 7 $\beta/\alpha, \gamma/\beta$

Provability (syntactic consequence) in CPL

$$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$$

1. α

premise

2. β

premise

3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$

Ax. 7 $\beta/\alpha, \gamma/\beta$

α β β

4. $\beta \rightarrow (\alpha \rightarrow \alpha)$

Ax. 1 $\alpha/\beta, \beta/\alpha$

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α

premise

2. β

premise

3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$

Ax. 7 $\beta/\alpha, \gamma/\beta$

4. $\beta \rightarrow (\alpha \rightarrow \alpha)$

Ax. 1 $\alpha/\beta, \beta/\alpha$

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α

2. β

3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$

4. $\beta \rightarrow (\alpha \rightarrow \alpha)$

5. $\alpha \rightarrow \alpha$

premise

premise

Ax. 7 $\beta/\alpha, \gamma/\beta$

Ax. 1 $\alpha/\beta, \beta/\alpha$

MP 4, 2

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α

premise

2. β

premise

3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$

Ax. 7 $\beta/\alpha, \gamma/\beta$

4. $\beta \rightarrow (\alpha \rightarrow \alpha)$

Ax. 1 $\alpha/\beta, \beta/\alpha$

5. $\alpha \rightarrow \alpha$

MP 4, 2

Provability (syntactic consequence) in CPL

$$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$$

1. α

premise

2. β

premise

3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$

Ax. 7 $\beta/\alpha, \gamma/\beta$

4. $\beta \rightarrow (\alpha \rightarrow \alpha)$

Ax. 1 $\alpha/\beta, \beta/\alpha$

5. $\alpha \rightarrow \alpha$

MP 4, 2

6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$

MP 3, 5

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α

premise

2. β

premise

3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$

Ax. 7 $\beta/\alpha, \gamma/\beta$

4. $\beta \rightarrow (\alpha \rightarrow \alpha)$

Ax. 1 $\alpha/\beta, \beta/\alpha$

5. $\alpha \rightarrow \alpha$

MP 4, 2

6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$

MP 3, 5

7. $(\overset{\alpha}{\alpha} \rightarrow (\overset{\beta}{\beta} \rightarrow \overset{\gamma}{\beta})) \rightarrow (\overset{\beta}{\beta} \rightarrow (\overset{\alpha}{\alpha} \rightarrow \overset{\gamma}{\beta}))$

Ax. 4 γ/β

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α

premise

2. β

premise

3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$

Ax. 7 $\beta/\alpha, \gamma/\beta$

4. $\beta \rightarrow (\alpha \rightarrow \alpha)$

Ax. 1 $\alpha/\beta, \beta/\alpha$

5. $\alpha \rightarrow \alpha$

MP 4, 2

6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$

MP 3, 5

7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$

Ax. 4 γ/β

8. $\alpha \rightarrow (\beta \rightarrow \beta)$

Ax. 1

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α	premise
2. β	premise
3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$	Ax. 7 $\beta/\alpha, \gamma/\beta$
4. $\beta \rightarrow (\alpha \rightarrow \alpha)$	Ax. 1 $\alpha/\beta, \beta/\alpha$
5. $\alpha \rightarrow \alpha$	MP 4, 2
6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$	MP 3, 5
7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$	Ax. 4 γ/β
8. $\alpha \rightarrow (\beta \rightarrow \beta)$	Ax. 1
9. $\beta \rightarrow (\alpha \rightarrow \beta)$	MP 7, 8

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α

premise

2. β

premise

3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$

Ax. 7 $\beta/\alpha, \gamma/\beta$

4. $\beta \rightarrow (\alpha \rightarrow \alpha)$

Ax. 1 $\alpha/\beta, \beta/\alpha$

5. $\alpha \rightarrow \alpha$

MP 4, 2

6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$

MP 3, 5

7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$

Ax. 4 γ/β

8. $\alpha \rightarrow (\beta \rightarrow \beta)$

Ax. 1

9. $\beta \rightarrow (\alpha \rightarrow \beta)$

MP 7, 8

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α	premise
2. β	premise
3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$	Ax. 7 $\beta/\alpha, \gamma/\beta$
4. $\beta \rightarrow (\alpha \rightarrow \alpha)$	Ax. 1 $\alpha/\beta, \beta/\alpha$
5. $\alpha \rightarrow \alpha$	MP 4, 2
6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$	MP 3, 5
7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$	Ax. 4 γ/β
8. $\alpha \rightarrow (\beta \rightarrow \beta)$	Ax. 1
9. $\beta \rightarrow (\alpha \rightarrow \beta)$	MP 7, 8
10. $\alpha \rightarrow \beta$	MP 9, 2

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α	premise
2. β	premise
3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$	Ax. 7 $\beta/\alpha, \gamma/\beta$
4. $\beta \rightarrow (\alpha \rightarrow \alpha)$	Ax. 1 $\alpha/\beta, \beta/\alpha$
5. $\alpha \rightarrow \alpha$	MP 4, 2
6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$	MP 3, 5
7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$	Ax. 4 γ/β
8. $\alpha \rightarrow (\beta \rightarrow \beta)$	Ax. 1
9. $\beta \rightarrow (\alpha \rightarrow \beta)$	MP 7, 8
10. $\alpha \rightarrow \beta$	MP 9, 2

Provability (syntactic consequence) in CPL

$$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$$

1. α	premise
2. β	premise
3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$	Ax. 7 $\beta/\alpha, \gamma/\beta$
4. $\beta \rightarrow (\alpha \rightarrow \alpha)$	Ax. 1 $\alpha/\beta, \beta/\alpha$
5. $\alpha \rightarrow \alpha$	MP 4, 2
6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$	MP 3, 5
7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$	Ax. 4 γ/β
8. $\alpha \rightarrow (\beta \rightarrow \beta)$	Ax. 1
9. $\beta \rightarrow (\alpha \rightarrow \beta)$	MP 7, 8
10. $\alpha \rightarrow \beta$	MP 9, 2

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α	premise
2. β	premise
3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$	Ax. 7 $\beta/\alpha, \gamma/\beta$
4. $\beta \rightarrow (\alpha \rightarrow \alpha)$	Ax. 1 $\alpha/\beta, \beta/\alpha$
5. $\alpha \rightarrow \alpha$	MP 4, 2
6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$	MP 3, 5
7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$	Ax. 4 γ/β
8. $\alpha \rightarrow (\beta \rightarrow \beta)$	Ax. 1
9. $\beta \rightarrow (\alpha \rightarrow \beta)$	MP 7, 8
10. $\alpha \rightarrow \beta$	MP 9, 2
11. $\alpha \rightarrow (\alpha \wedge \beta)$	MP 6, 10

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α	premise
2. β	premise
3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$	Ax. 7 $\beta/\alpha, \gamma/\beta$
4. $\beta \rightarrow (\alpha \rightarrow \alpha)$	Ax. 1 $\alpha/\beta, \beta/\alpha$
5. $\alpha \rightarrow \alpha$	MP 4, 2
6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$	MP 3, 5
7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$	Ax. 4 γ/β
8. $\alpha \rightarrow (\beta \rightarrow \beta)$	Ax. 1
9. $\beta \rightarrow (\alpha \rightarrow \beta)$	MP 7, 8
10. $\alpha \rightarrow \beta$	MP 9, 2
11. $\alpha \rightarrow (\alpha \wedge \beta)$	MP 6, 10

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α	premise
2. β	premise
3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$	Ax. 7 $\beta/\alpha, \gamma/\beta$
4. $\beta \rightarrow (\alpha \rightarrow \alpha)$	Ax. 1 $\alpha/\beta, \beta/\alpha$
5. $\alpha \rightarrow \alpha$	MP 4, 2
6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$	MP 3, 5
7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$	Ax. 4 γ/β
8. $\alpha \rightarrow (\beta \rightarrow \beta)$	Ax. 1
9. $\beta \rightarrow (\alpha \rightarrow \beta)$	MP 7, 8
10. $\alpha \rightarrow \beta$	MP 9, 2
11. $\alpha \rightarrow (\alpha \wedge \beta)$	MP 6, 10
12. $\alpha \wedge \beta$	MP 11, 1

Provability (syntactic consequence) in CPL

$\{\alpha, \beta\} \vdash_p \alpha \wedge \beta$

1. α	premise
2. β	premise
3. $(\alpha \rightarrow \alpha) \rightarrow ((\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta)))$	Ax. 7 $\beta/\alpha, \gamma/\beta$
4. $\beta \rightarrow (\alpha \rightarrow \alpha)$	Ax. 1 $\alpha/\beta, \beta/\alpha$
5. $\alpha \rightarrow \alpha$	MP 4, 2
6. $(\alpha \rightarrow \beta) \rightarrow (\alpha \rightarrow (\alpha \wedge \beta))$	MP 3, 5
7. $(\alpha \rightarrow (\beta \rightarrow \beta)) \rightarrow (\beta \rightarrow (\alpha \rightarrow \beta))$	Ax. 4 γ/β
8. $\alpha \rightarrow (\beta \rightarrow \beta)$	Ax. 1
9. $\beta \rightarrow (\alpha \rightarrow \beta)$	MP 7, 8
10. $\alpha \rightarrow \beta$	MP 9, 2
11. $\alpha \rightarrow (\alpha \wedge \beta)$	MP 6, 10
12. $\alpha \wedge \beta$	MP 11, 1

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Is there a complete match, i.e.,

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Is there a complete match, i.e.,

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Is there a complete match, i.e.,

$$\{\text{Tautology}\} = \{\text{Theorems}\}$$

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Is there a complete match, i.e.,

$\{\text{Tautology}\} = \{\text{Theorems}\}$

$\Phi \models \alpha \quad \text{iff} \quad \Phi \vdash_p \alpha \quad ?$

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Is there a complete match, i.e.,

$\{\text{Tautology}\} = \{\text{Theorems}\}$

$\Phi \models \alpha \quad \text{iff} \quad \Phi \vdash_p \alpha \quad ?$

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Is there a complete match, i.e.,

$\{\text{Tautology}\} = \{\text{Theorems}\}$

$\Phi \models \alpha \quad \text{iff} \quad \Phi \vdash_p \alpha \quad ?$

In CPL, the answer is YES.

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Is there a complete match, i.e.,

$\{\text{Tautology}\} = \{\text{Theorems}\}$

$\Phi \models \alpha \quad \text{iff} \quad \Phi \vdash_p \alpha \quad ?$

In CPL, the answer is YES.

Big questions

How are semantic concept (tautology, semantic consequence) related to syntactic concepts (theorems, provability)?

Is there a complete match, i.e.,

$\{\text{Tautology}\} = \{\text{Theorems}\}$

$\Phi \models \alpha \quad \text{iff} \quad \Phi \vdash_p \alpha \quad ?$

In CPL, the answer is YES.

Completeness of CPL

COMPLETENESS THEOREM

For an arbitrary wff α of CPL:

α is a theorem of CPL iff α is a tautology of CPL.

Truth vs proof

Truth vs proof

That was CPL. Could we do it more abstractly? E.g., in CPL we assumed two true-values, truth and falsehood. Could we *define* truth, and then ask what the relation between truth and proof is?

Truth vs proof

That was CPL. Could we do it more abstractly? E.g., in CPL we assumed two true-values, truth and falsehood. Could we *define* truth, and then ask what the relation between truth and proof is?

Yes, Alfred Tarski's project and his (restricted) success: pick a formalized language and define truth for sentences in this language in such a way that the Liar's paradox is not derivable.

Truth vs proof

That was CPL. Could we do it more abstractly? E.g., in CPL we assumed two true-values, truth and falsehood. Could we *define* truth, and then ask what the relation between truth and proof is?

Yes, Alfred Tarski's project and his (restricted) success: pick a formalized language and define truth for sentences in this language in such a way that the Liar's paradox is not derivable.

Having a formalized language with the concept of truth for that language, consider axiom systems in that language. Aim for the best system. And ask: is there an axiom system that yields as theorems true and only true sentences of that language?

Truth vs proof

That was CPL. Could we do it more abstractly? E.g., in CPL we assumed two true-values, truth and falsehood. Could we *define* truth, and then ask what the relation between truth and proof is?

Yes, Alfred Tarski's project and his (restricted) success: pick a formalized language and define truth for sentences in this language in such a way that the Liar's paradox is not derivable.

Having a formalized language with the concept of truth for that language, consider axiom systems in that language. Aim for the best system. And ask: is there an axiom system that yields as theorems true and only true sentences of that language?

Tarski / Goedel result: for any sufficiently rich language, the set of true sentences is larger than the set of provable sentences (in any axiom system).

Truth vs proof

Truth vs. proof is an exam topic. Literature:

- (1) Some logic textbook to explain what a formal language is and the machinery of Tarski's truth definition.
- (2) Initial paragraphs of Tarski's classic "Truth in formalized languages" that explains the philosophy of his project (the Liar's paradox and some history of attempts to define truth)
- (3) Tarski's paper "Truth and proof".